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# *Thread and Synchronization*

## *Task Model (Module 18)*

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# Why Talk About This Subject

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## □ A thread of program execution

- ❖ How a program start and end its execution
- ❖ waiting for an event or a resource, delay a period, etc.

## □ For concurrent operations → multiple threads of program execution

## □ How can we make this happen?

- ❖ support for program execution
- ❖ sharing of resources
- ❖ scheduling
- ❖ communication between threads



# Thread and Process

## □ Process:

- ❖ an entity to which system resources (CPU time, memory, etc.) are allocated
- ❖ an address space with 1 or more threads executing within that address space, and the required system resources for those threads

## □ Thread:

- ❖ a sequence of control within a process and shares the resources in that process

## □ Lightweight process (LWP):

- ❖ LWP may share resources: address space, open files, ...
- ❖ clone or fork – share or not share address space, file descriptor, etc.
- ❖ In Linux kernel, threads are implemented as standard processes (LWP) that shares certain resources with other processes, and there is no special scheduling semantics or data structures to represent threads



# Why Threads

## □ Advantages:

- ❖ the overhead for creating a thread is significantly less than that for creating a process
- ❖ multitasking, i.e., one process serves multiple clients
- ❖ switching between threads requires the OS to do much less work than switching between processes

## □ Drawbacks:

- ❖ not as widely available as the process features
- ❖ writing multithreaded programs require more careful thought
- ❖ more difficult to debug than single threaded programs
- ❖ for single processor machines, creating several threads in a program may not necessarily produce an increase in performance (only so many CPU cycles to be had)



# POSIX Thread (pthread)

## ❑ IEEE's POSIX Threads Model:

- ❖ programming models for threads in a UNIX platform
- ❖ pthreads are included in the international standards

## ❑ pthreads programming model:

- ❖ creation of threads
- ❖ managing thread execution
- ❖ managing the shared resources of the process

## ❑ main thread:

- ❖ initial thread created when main() is invoked
- ❖ has the ability to create daughter threads
- ❖ if the main thread returns, the process terminates even if there are running threads in that process
- ❖ to explicitly avoid terminating the entire process, use pthread\_exit()



# Linux task\_struct

```
struct task_struct {                                /* Linux/include/linux/sched.h */
    volatile long state; /* -1 unrunnable, 0 runnable, >0 stopped */
    void *stack;
    atomic_t usage;
    unsigned int flags; /* per process flags, defined below */
    unsigned int ptrace;

    int lock_depth; /* BKL (big kernel lock) lock depth */

    int prio, static_prio, normal_prio;
    unsigned int rt_priority;
    const struct sched_class *sched_class;
    .....
    struct mm_struct *mm, *active_mm;
    struct thread_struct thread; /* CPU-specific state of this task */
    struct fs_struct *fs; /* filesystem information */
    struct files_struct *files; /* open file information */
};
```



# Process -- *task\_struct* data structure

## ❑ state: process state

- ❖ TASK\_RUNNING: executing
- ❖ TASK\_INTERRUPTABLE: suspended (sleeping)
- ❖ TASK\_UNINTERRUPTABLE: (no process of signals)
- ❖ TASK\_STOPPED (stopped by SIGSTOP)
- ❖ TASK\_TRACED (being monitored by other processes such as debuggers)
- ❖ EXIT\_ZOMBIE (terminated before waiting for parent)
- ❖ EXIT\_DEAD

## ❑ thread\_info: low-level information for the process

## ❑ mm: pointers to memory area descriptors

## ❑ tty: tty associated with the process

## ❑ fs: current directory

## ❑ files: pointers to file descriptors

## ❑ signal: signals received .....



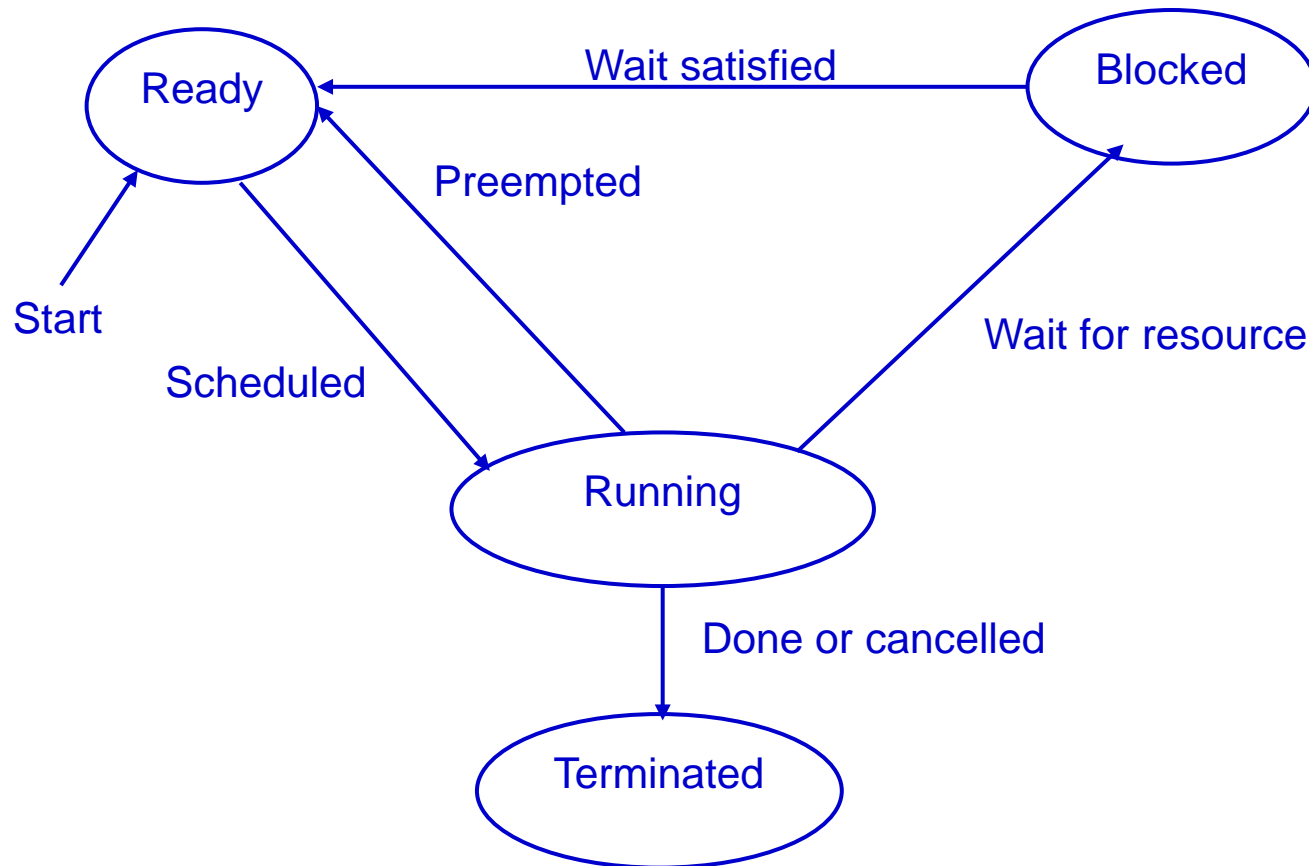
# Linux Processor State

```
/* This is the TSS (task State Segment) defined by the hardware and saved in stack. */
struct x86_hw_tss {
    unsigned short    back_link, __blh;
    unsigned long     sp0;
    unsigned short    ss0, __ss0h;
    unsigned long     sp1;
    unsigned short    ss1, __ss1h; /* ss1 caches MSR_IA32_SYSENTER_CS: */
    unsigned long     sp2;
    unsigned short    ss2, __ss2h;
    unsigned long     __cr3;
    unsigned long     ip;
    unsigned long     flags;
    unsigned long     ax;
    unsigned long     cx;
    unsigned long     dx;
    unsigned long     bx;
/* For ARM, Linux/arch/arm/include/asm/thread_info.h.,
```

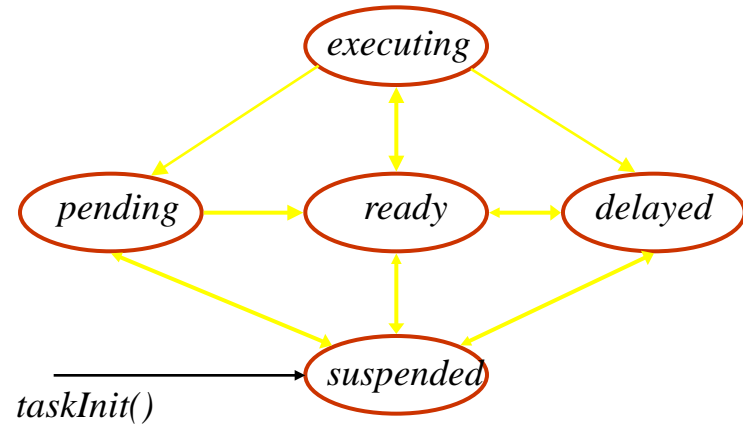
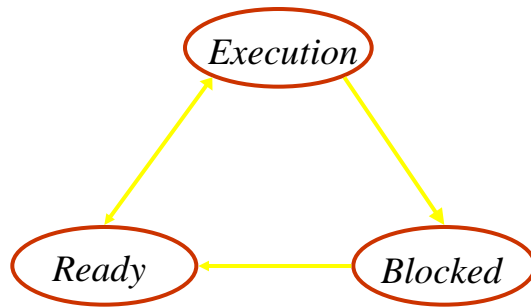




# Linux Thread State Transition



# Task Management in vxWorks



## □ Task structure in task control block –

- ❖ priority(initial and inherited), stack frame, task current state,
- ❖ entry point, processor states (program counter, registers)
- ❖ callback function (hook) pointers for OS events
- ❖ spare variables



# VxWorks Task States

```
typedef struct windTcb          /* WIND_TCB - task control block */
{
    char *          name;      /* 0x34: pointer to task name */
    UINT           status;     /* 0x3c: status of task */
    UINT           priority;   /* 0x40: task's current priority */
    UINT           priNormal;  /* 0x44: task's normal priority */
    UINT           priMutexCnt; /* 0x48: nested priority mutex owned */
    UINT           lockCnt;    /* 0x50: preemption lock count */
    FUNCPTR        entry;     /* 0x74: entry point of task */
    char *          pStackBase; /* 0x78: points to bottom of stack */
    char *          pStackLimit; /* 0x7c: points to stack limit */
    char *          pStackEnd; /* 0x80: points to init stack limit */
#ifdef CPU_FAMILY==I80X86 /* function declarations */
    EXC_INFO        exclInfo;  /* 0x118: exception info */
    REG_SET         regs;      /* 0x12c: register set */
    DBG_INFO_NEW    dbgInfo0;  /* 0x154: debug info */
#endif /* CPU_FAMILY==I80X86 */
};
```

